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## ***HEC POSIX I/O API Extensions***

***Rob Ross***

***Mathematics and Computer Science Division***

***Argonne National Laboratory***

***[rross@mcs.anl.gov](mailto:rross@mcs.anl.gov)***

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## POSIX Introduction

- POSIX is the IEEE Portable Operating System Interface for Computing Environments.
- “POSIX defines a standard way for an application program to obtain basic services from the operating system”
- POSIX was created when a single computer owned its own file system.
  - Network file systems like NFS chose not to implement strict POSIX semantics in all cases (e.g., lazy access time propagation)
  - Heavily shared files (e.g., from clusters) can be very expensive for file systems that provide POSIX semantics, or have undefined contents for file systems that bend the rules
- The Open Group (<http://www.opengroup.org/>) is responsible for the specification and any subsequent extensions.

## APIs for HEC I/O

- POSIX IO APIs (open, close, read, write, stat) have semantics that can make it hard to achieve high performance when large clusters of machines access shared storage.
- A working group of HEC users is drafting some proposed API additions for POSIX that will provide standard ways to achieve higher performance.
- Primary approach is either to relax semantics that can be expensive, or to provide more information to inform the storage system about access patterns.
- The goal is to create a standard way to provide high performance and good semantics
- Three components:
  - **Good concepts** - building blocks for more effective I/O systems
  - **API definition and standardization** - well-defined and capable interfaces to use these ideas agreed upon by the community
  - **Implementations** - early prototypes to show viability, adoption in OSeS and file systems to provide availability

## Contributors

- Lee Ward - Sandia National Lab
- Bill Lowe, Tyce McLarty – Lawrence Livermore National Lab
- Gary Grider, James Nunez – Los Alamos National Lab
- Rob Ross, Rajeev Thakur, William Gropp, Murali Vilayannur - Argonne National Lab
- Roger Haskin – IBM
- Brent Welch, Marc Unangst - Panasas
- Garth Gibson - Carnegie Mellon University/Panasas
- Alok Choudhary – Northwestern University
- Tom Ruwart - University of Minnesota/IO Performance
- Harriet Coverston - Sun Microsystems
- Others ...

## Current HEC POSIX Enhancement Areas

- Current “first string”:
  - Metadata (lazy attributes, aggregation)
    - *statlite()* and friends
    - *readdirplus()* and friends
  - Coherence – (last writer wins and other such things can be optional)
    - *O\_LAZY*, *lazyio\_propagate()*, *lazyio\_synchronize()*
  - Efficient name resolution and file open (group file opens)
    - *openg()*, *openfh()*
  - Flexible (if not concise) description of I/O operations
    - *readx()*, *writex()*
- Group locks, ACLs, QoS, and portable hinting are being investigated as well, but I will focus on the first string.

Please don't ask me about the specifics of the semantics of these calls during this talk! Instead ask me off-line, or join the mailing list (see web site on last slide) and engage the whole group.

## *statlite, fstatlite, lstatlite - Lazy Attributes*

### ■ Syntax

```
int statlite(const char *file_name, struct statlite *buf);  
int fstatlite(int filedes, struct statlite *buf);  
int lstatlite(const char *file_name, struct statlite *buf);
```

### ■ Description

- This family of stat calls, the lite family, is provided to allow for file I/O performance not to be compromised by frequent use of stat information lookup. Some information can be expensive to obtain when a file is busy.
- They all return a *stat* structure, which has all the normal fields from the stat family of calls but some of the fields (e.g., file size, modify time) are optionally not guaranteed to be correct.
- There is a litemask field that can be used to specify which of the optional fields you require to be completely correct values returned.
- **statlite** stats the file pointed to by *file\_name* and fills in *buf*.
- **fstatlite** is identical to **stat**, only the open file pointed to by *filedes* (as returned by **open**(2)) is statlited-ed in place of *file\_name*.

## *readdirplus, readdirlite - Aggregating Metadata Operations*

### ■ Syntax

```
struct dirent_plus *readdirplus(DIR *dirp);
```

```
int readdirplus_r(DIR *dirp, struct dirent_plus *entry, struct dirent_plus  
  **result);
```

```
struct dirent_lite *readdirlite(DIR *dirp);
```

```
int readdirlite_r(DIR *dirp, struct dirent_lite *entry, struct dirent_lite  
  **result);
```

### ■ Description

- This family of calls is provided to all the file system to return file metadata as part of the directory read process. This as a side-effect aggregates many stat operations together.
- **readdirplus**(2) and **readdirplus\_r**(2) return a directory entry plus **lstat**(2) results (like the NFSv3 REaddirplus command)
- **readdirlite**(2) and **readdirlite\_r**(2) return a directory entry plus **lstatlite**(2) results

## *O\_LAZY, lazyio\_propagate, lazyio\_synchronize* - Coherence

### ■ Syntax

Specify `O_LAZY` in *flags* argument to `open(2)`

```
int lazyio_propagate(int fd, off_t offset, size_t count);
```

```
int lazyio_synchronize(int fd, off_t offset, size_t count);
```

### ■ Description

- Requests lazy I/O data integrity. Allows network filesystem to relax data coherency requirements to improve performance for shared-write file. This is a hint only: if filesystem does not support lazy I/O integrity, does not have to do anything differently.
- Writes may not be visible to other processes or clients until `lazyio_propagate(2)`, `fsync(2)`, or `close(2)` is called
- Reads may come from local cache (ignoring changes to file on backing storage) until `lazyio_synchronize(2)` is called
- Does not provide synchronization across processes or nodes – program must use external synchronization (e.g., pthreads, MPI, etc.) to coordinate actions.

## *openg, openfh - Name Space Traversal and Collective File Open*

### ■ Syntax

```
int openg(char *path, int mode, fh_t *handle);  
int openfh(fh_t *fh);
```

### ■ Description

- The **openg**() function opens a file named by path according to mode (e.g., O\_RDWR). It returns an opaque file handle corresponding to a file descriptor. The intent is that the file handle can be transferred to cooperating processes and converted to a file descriptor with **openfh**()).
- The **openfh**() function shall create an open file descriptor that refers to the file represented by the *fh* argument. The file status flags and file access modes of the open file description shall be set according to those given in the accompanying **openg**()).
- The lifetime of the file handle is implementation specific. For example, it may not be valid once all open file descriptors derived from the handle with **openfh**() have been closed.

## *readx, writex - Efficient I/O Description*

### ■ Syntax

```
ssize_t readx(int fd, const struct iovec *iov, size_t iov_count, struct
    xtvec *xtv, size_t xtv_count);
ssize_t writex(int fd, const struct iovec *iov, size_t iov_count, struct
    xtvec *xtv, size_t xtv_count);
struct xtvec { off_t xtv_off; /* Starting file offset */
    size_t xtv_len; /* Number of bytes */ };
```

### ■ Description

- Generalized file vector to memory vector transfer. Existing **readv()**, **writev()** specify a memory vector and do serial IO. The new **readx()**, **writex()** calls also read/write strided vectors to/from files, but regions may be processed in any order, and iov and xtv need not have the same number of elements.
- The **readx()** function reads xtv\_count blocks described by xtv from the file associated with the file descriptor fd into the iov\_count multiple buffers described by iov. The file offset is not changed.
- The **writex()** function writes at most xtv\_count blocks described by xtv into the file associated with the file descriptor fd from the iov\_count multiple buffers described by iov. The file offset is not changed.

## Current Status and Contact Information

- Ideas
  - Group has identified short-term and long-term goals for improvements
- Interface Specification
  - HEC Extensions working group formed with Open Group
  - Draft 0 specification nearing completion, includes calls discussed here
- Implementations
  - Prototypes of many calls have been implemented by ANL, UCSC, Sun, CFS/Cray, etc.
  - Source for many of these will be made available soon
- Go to the POSIX HPC I/O Extensions Web site for more information:  
[www.pdl.cmu.edu/posix/](http://www.pdl.cmu.edu/posix/)

